

A new post-quantum digital signature scheme

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SIAM-AG TU/e

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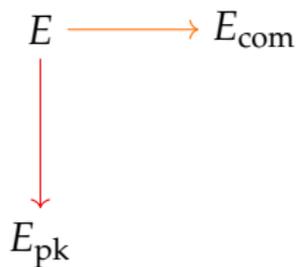
joint work (in progress) with Ross Bowden

Concept: Let's take the Q out of SQISign



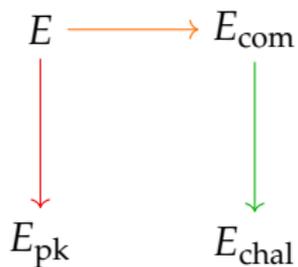
public, secret, commitment, challenge, response

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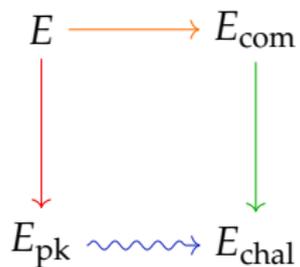
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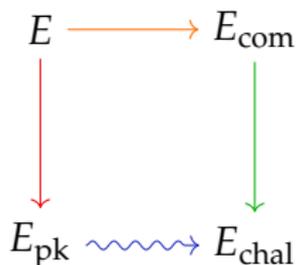
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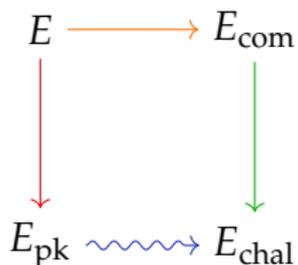
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- ▶ Bottleneck: response requires
 - ▶ Computing $\text{End}(E_{\text{chal}})$
 - ▶ KLPT

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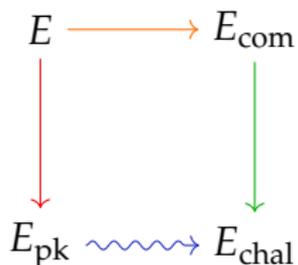


public, **secret**, **commitment**, **challenge**, response

- ▶ Bottleneck: **response** requires
 - ▶ Computing $\text{End}(E_{\text{chal}})$
 - ▶ KLPT
- ▶ Idea: replace
 - ▶ Computing $\text{End}(E_{\text{chal}}) \rightsquigarrow$ basis change on $T_2(E_{\text{pk}})$
 - ▶ KLPT \rightsquigarrow once-for-each- E_{pk} computation¹

¹Hopefully. Work in progress, remember?

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How?!

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How: Bruhat-Tits trees and isogeny graphs

Quaternion Algebra graphs

Nodes:

Maximal
orders / \sim

Edges:

Ideals of
norm 2 / \sim

Supersingular isogeny graphs

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Quotients of Bruhat-Tits trees

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Quotients of Bruhat-Tits trees

Nodes:

$$\left\{ \begin{pmatrix} 2^n & 0 \\ r & 2^m \end{pmatrix} : 0 \leq r < 2^m \right\} / \sim$$

Edges:

$$\begin{pmatrix} 2 & 0 \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 1 & 2 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 0 & 2 \end{pmatrix}$$

act on nodes via multiplication.

How: Bruhat-Tits trees and isogeny graphs

SIGs

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(Quotients of)
Bruhat-Tits trees

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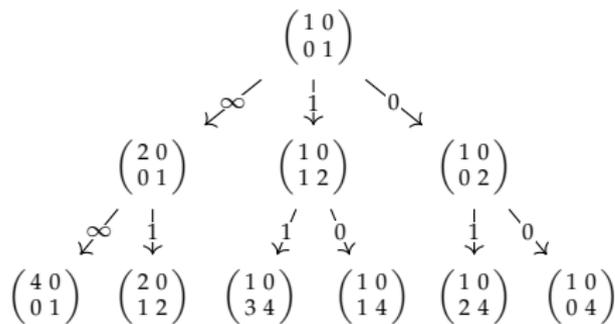
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1. Choose **root** of tree E_{pk} and **basis** $\{P, Q\}$ of $T_2(E_{pk})$.



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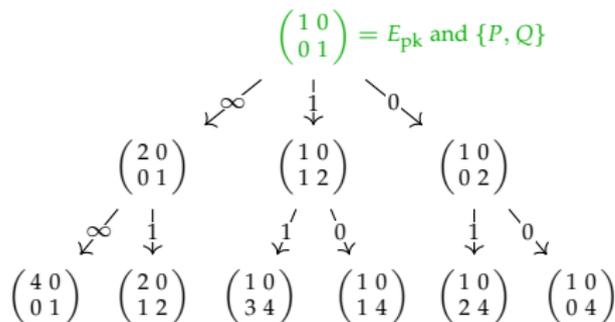
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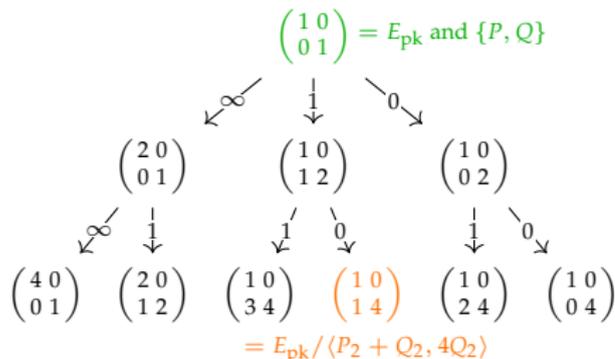
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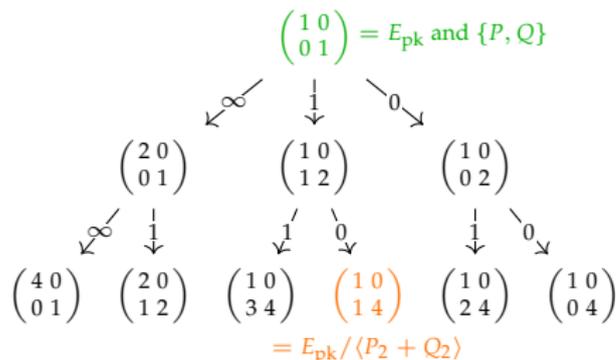
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2. Call P_d, Q_d canonical lifts of P, Q that generate $E_{pk}[2^d]$.
3. A **matrix node** M at depth d is

$$E_{pk} / \langle (P_d, Q_d) \cdot M \rangle.$$



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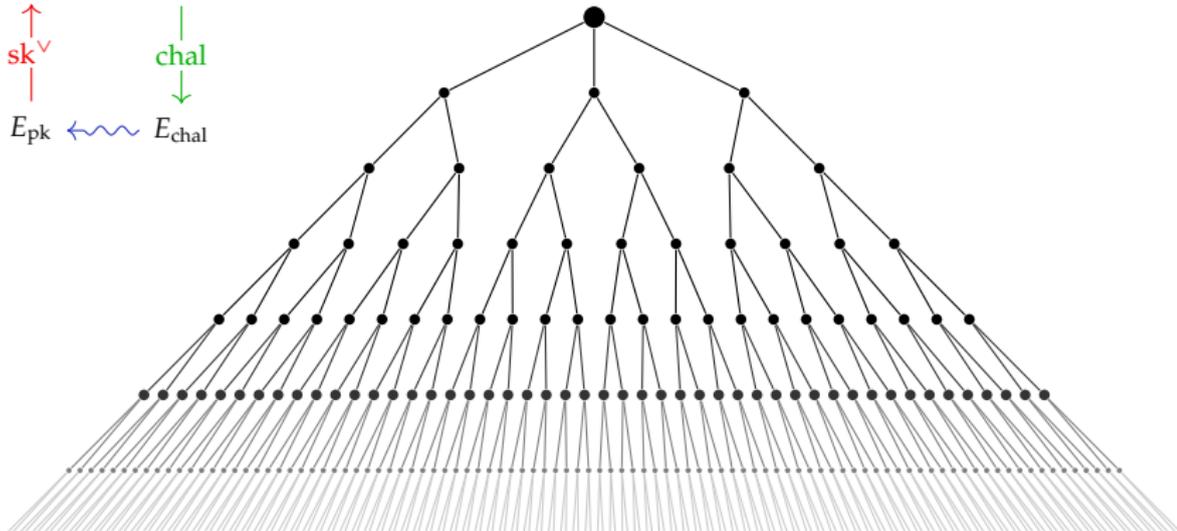
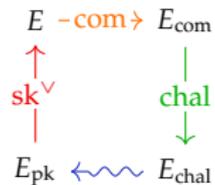
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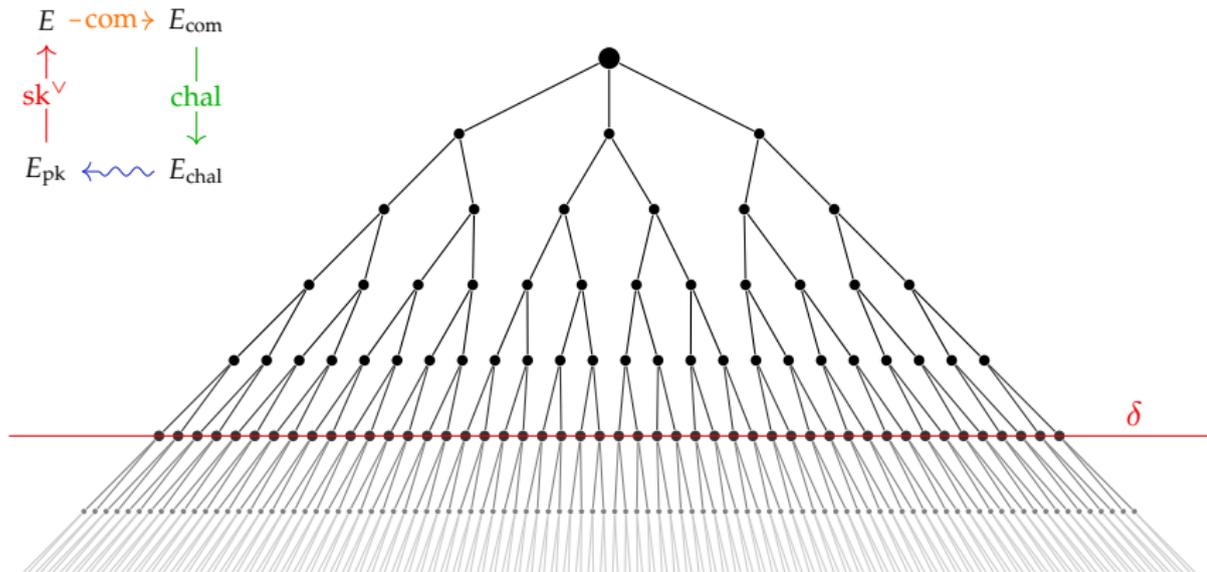
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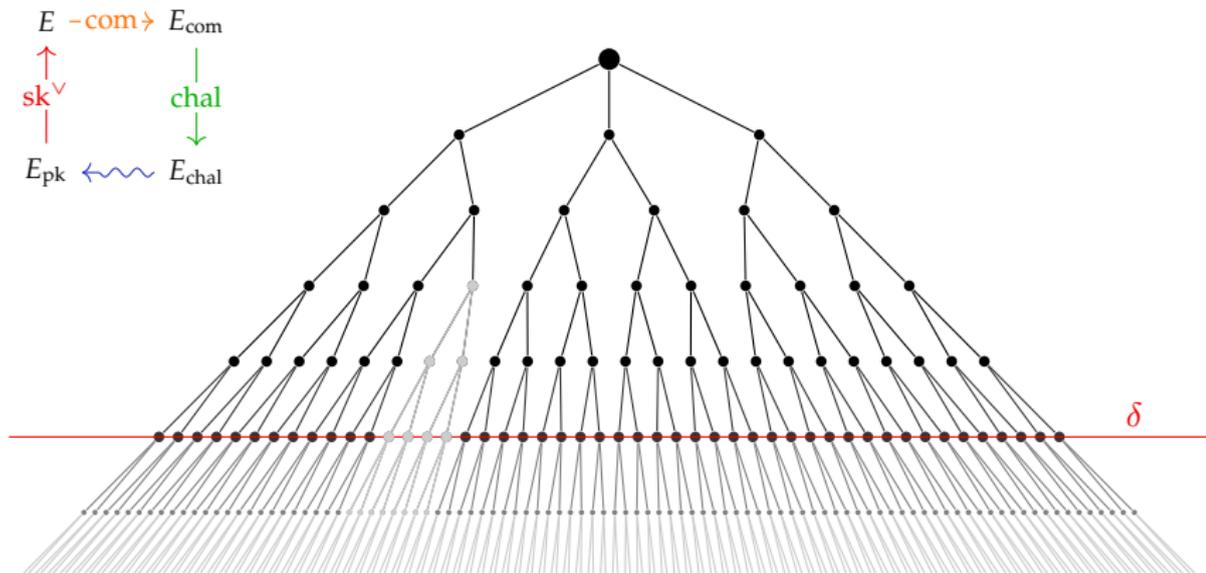
From KLPT to Bruhat-Tits quotients



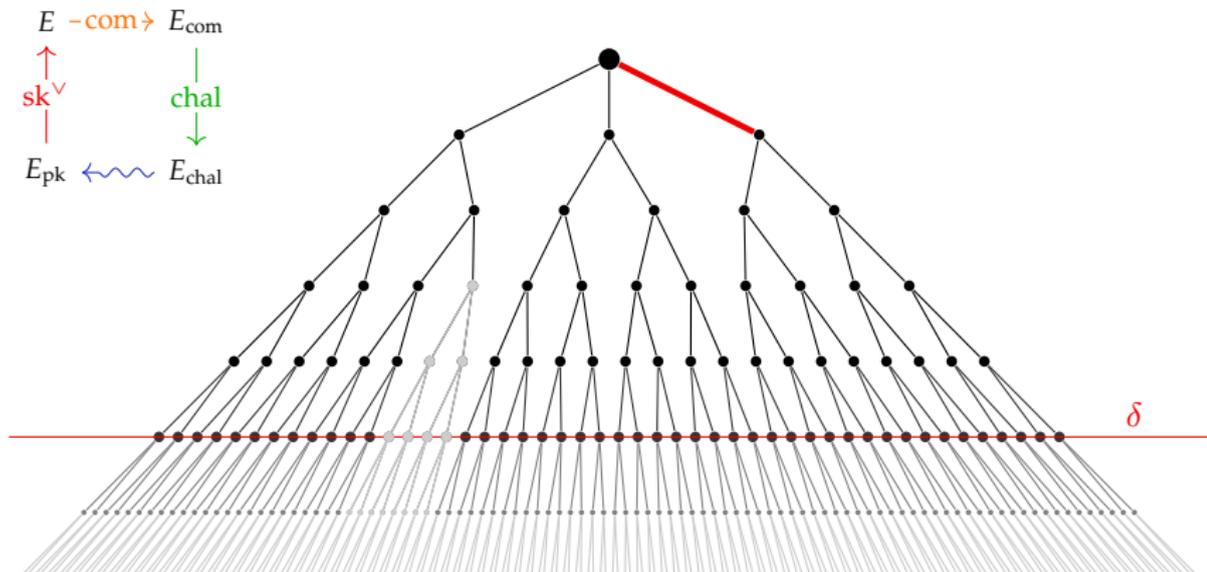
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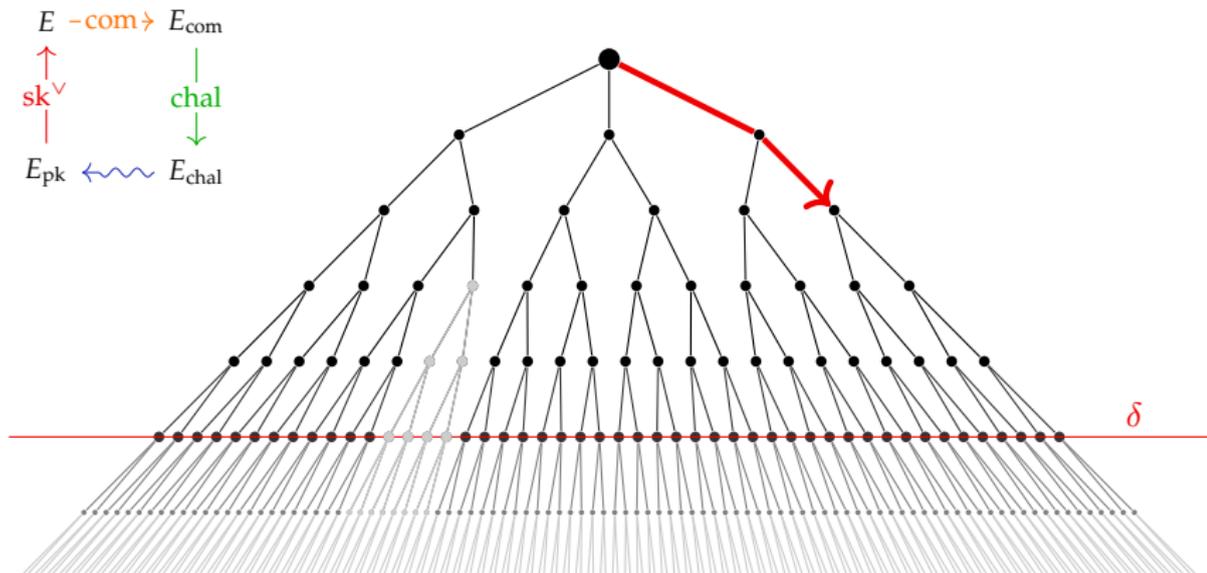
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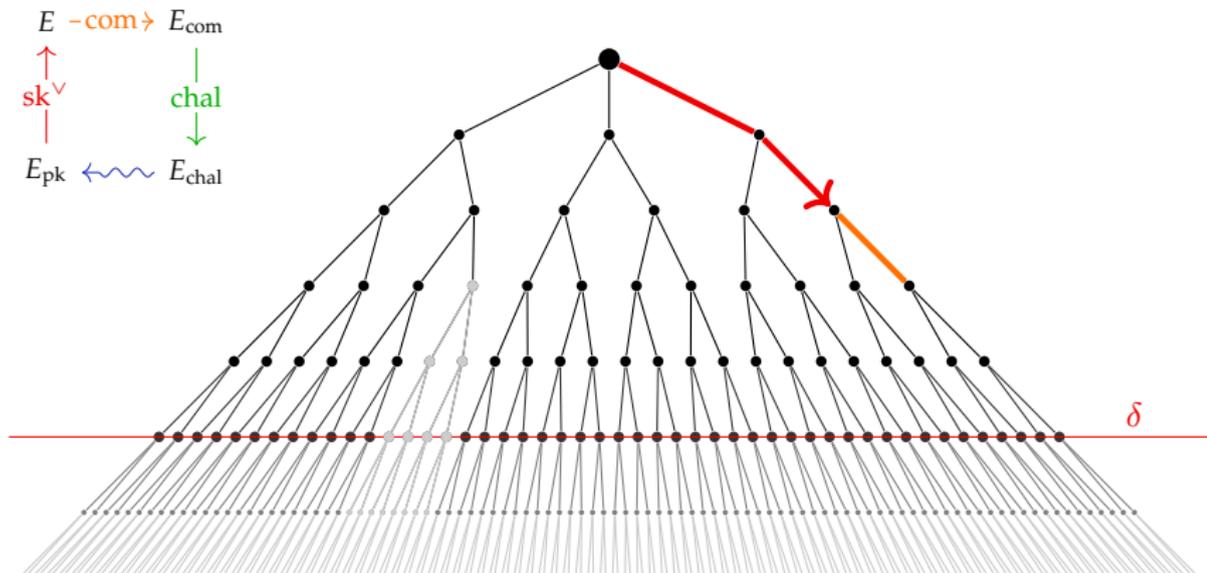
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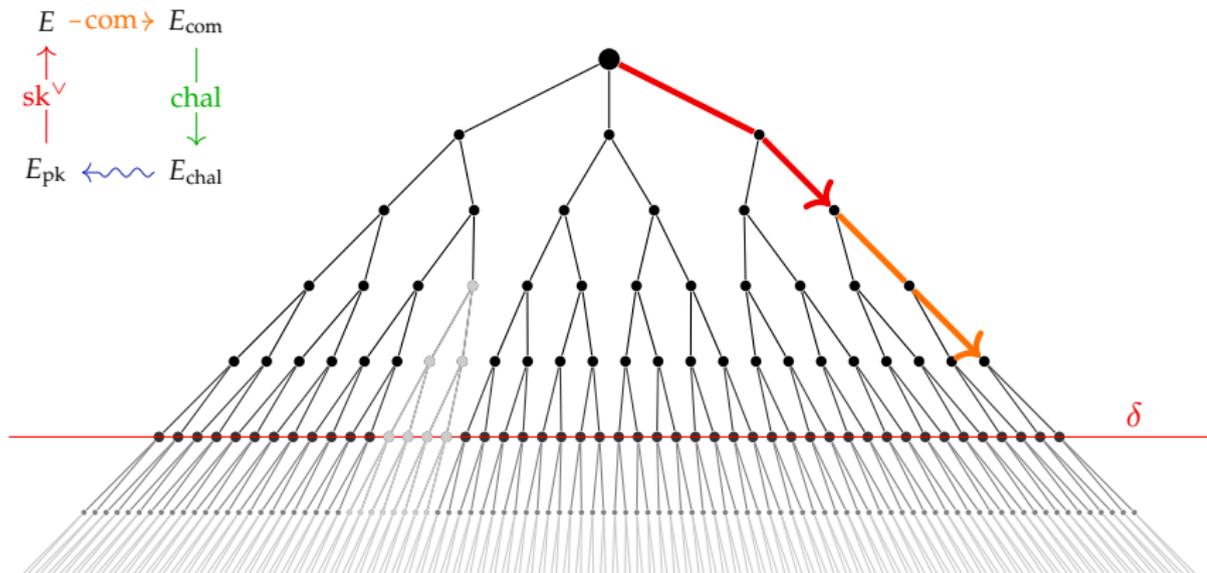
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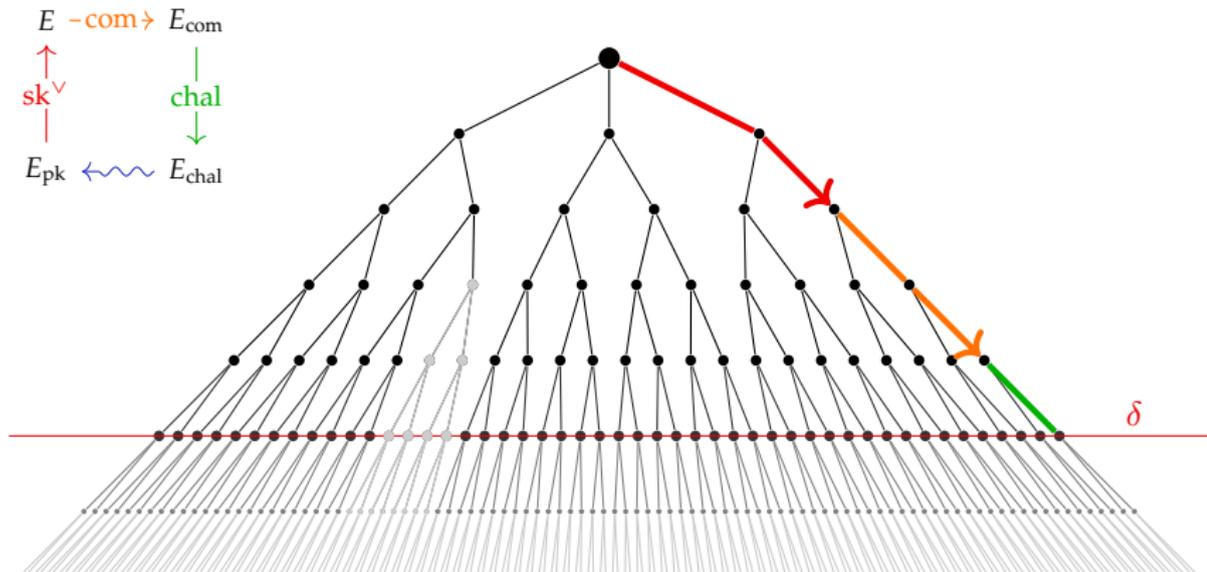
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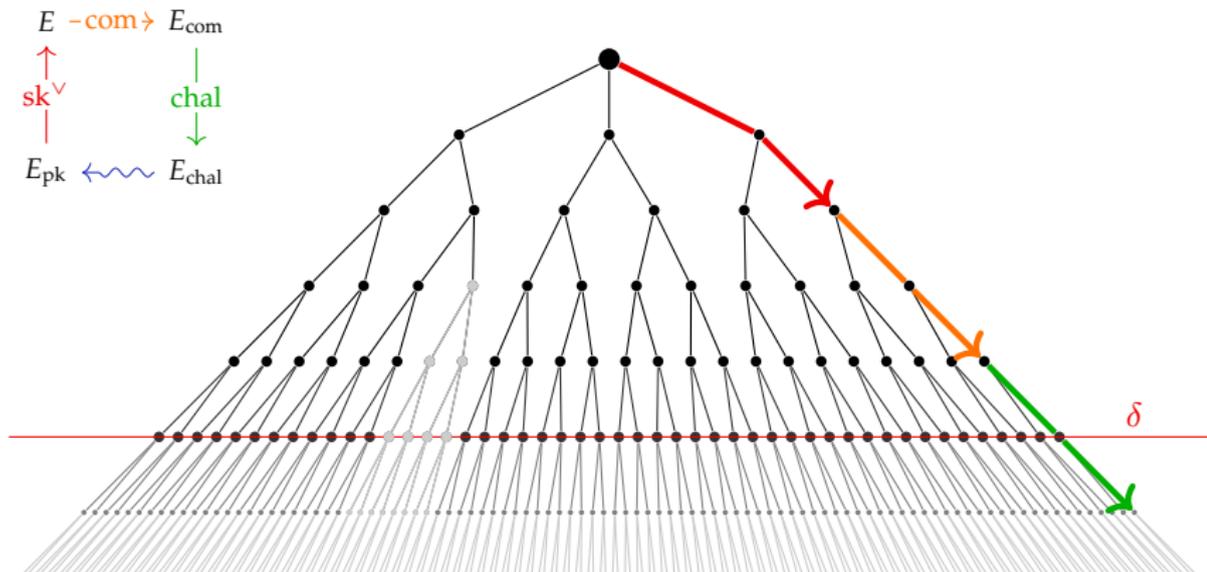
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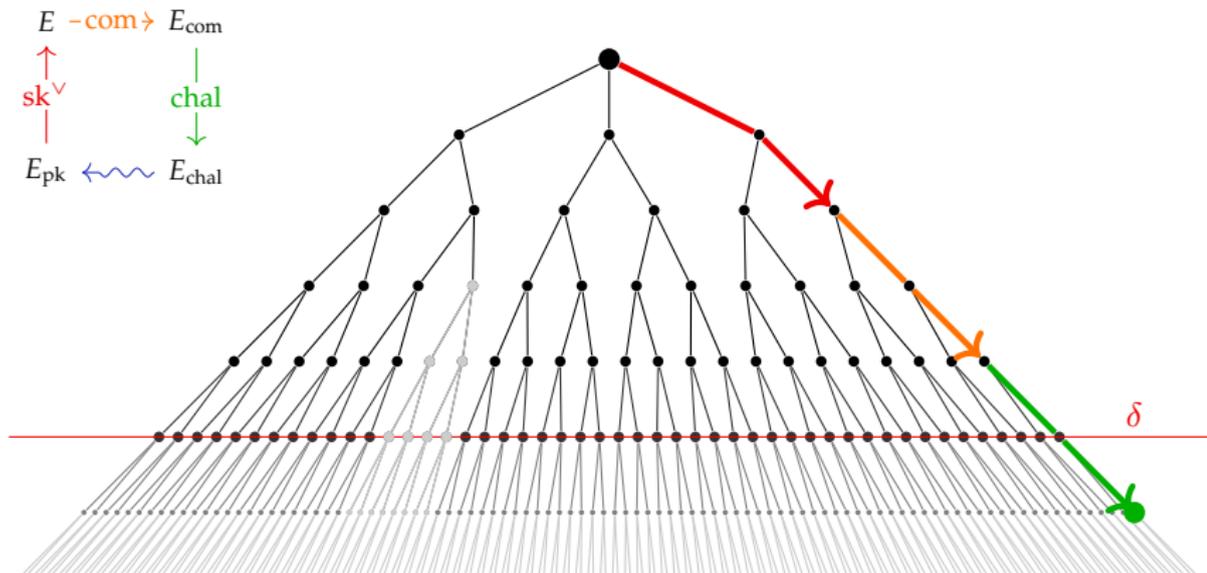
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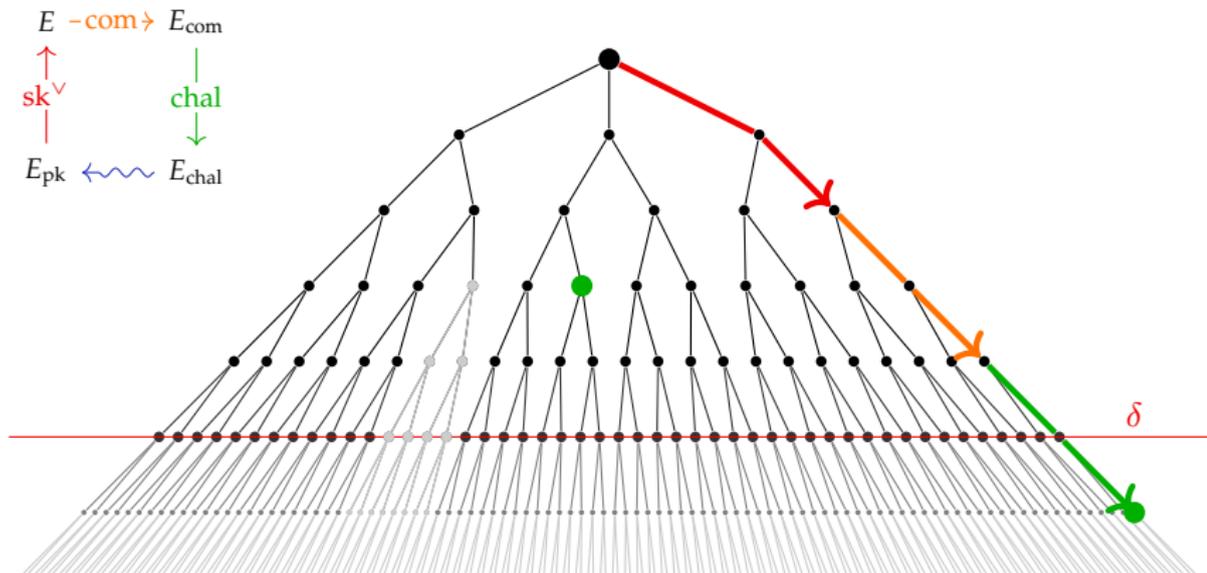
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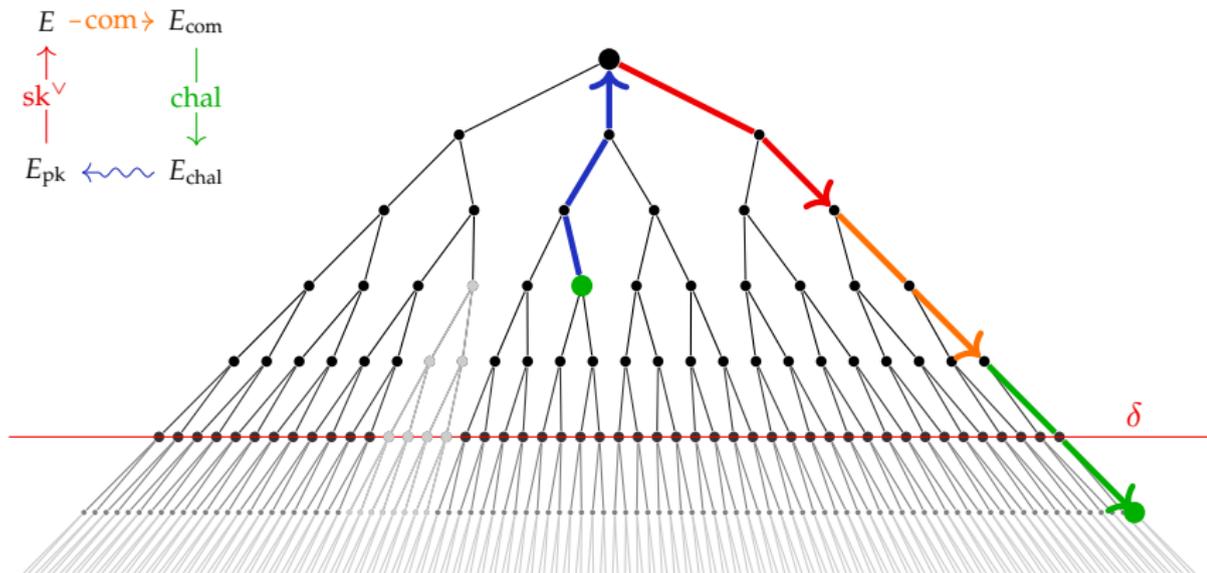
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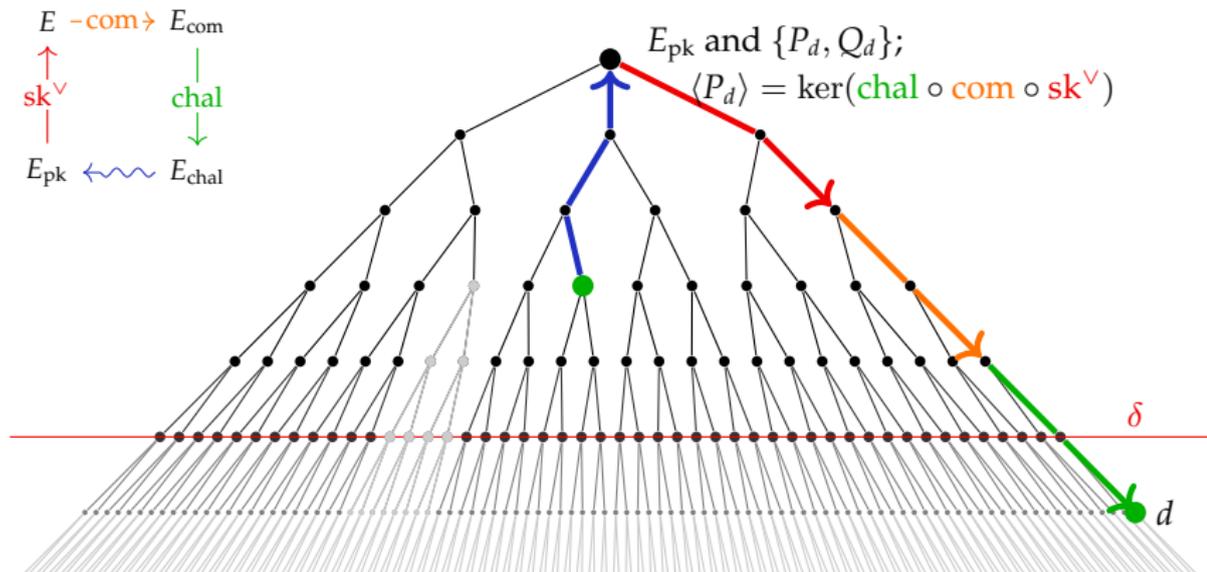
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The dream: A Bruhat-Tits identification protocol

Setup:

1. Starting at $j(E_0) = 1728$, compute a chain of 2-isogenies **sk**; publish codomain E_{pk} as public key.

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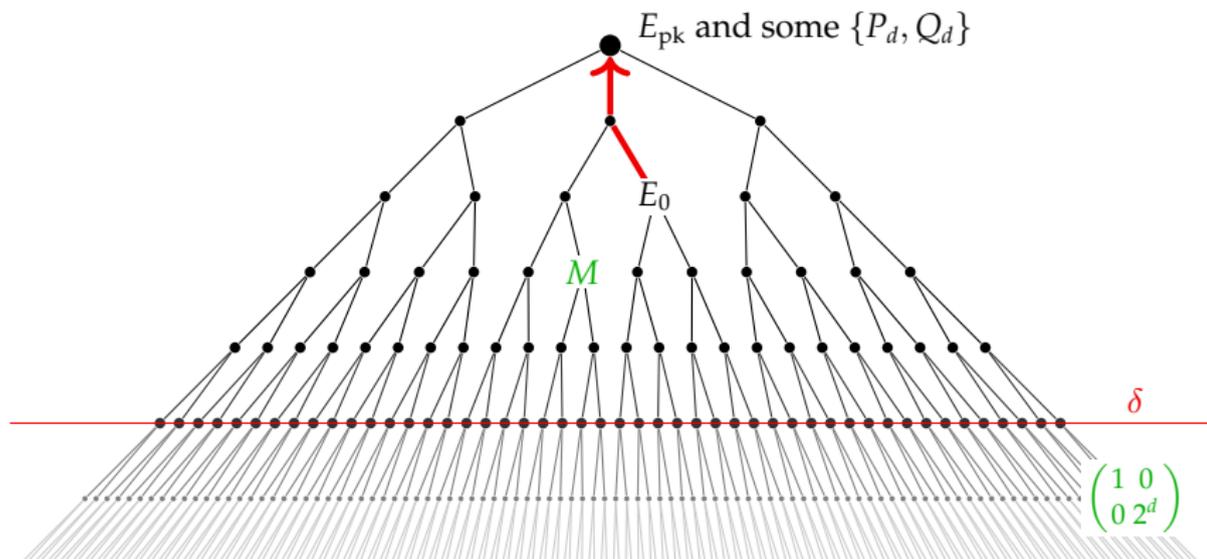
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3. Compute, for these choices, the matrix node M of depth $< \delta$ equivalent to $\begin{pmatrix} 1 & 0 \\ 0 & 2^d \end{pmatrix}$ in the quotient of the Bruhat-Tits tree with root E_{pk} and basis $\{P_d, Q_d\}$ of $T_2(E_{\text{pk}})$ at precision 2^d .

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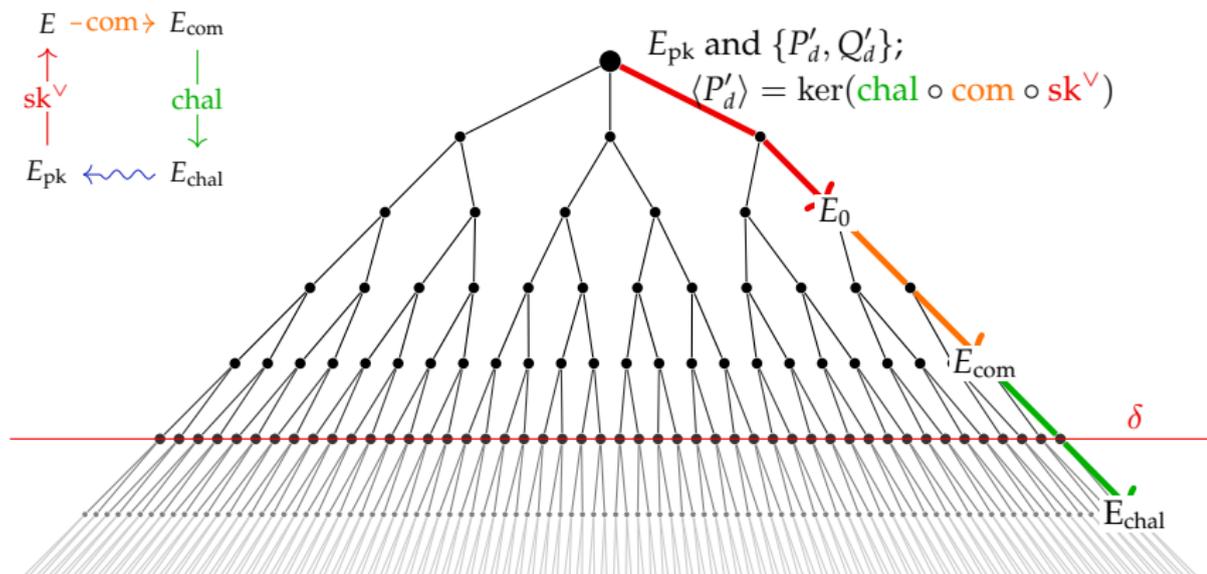
Response:

1. Compute the base change to basis $\{P'_d, Q'_d\}$ such that

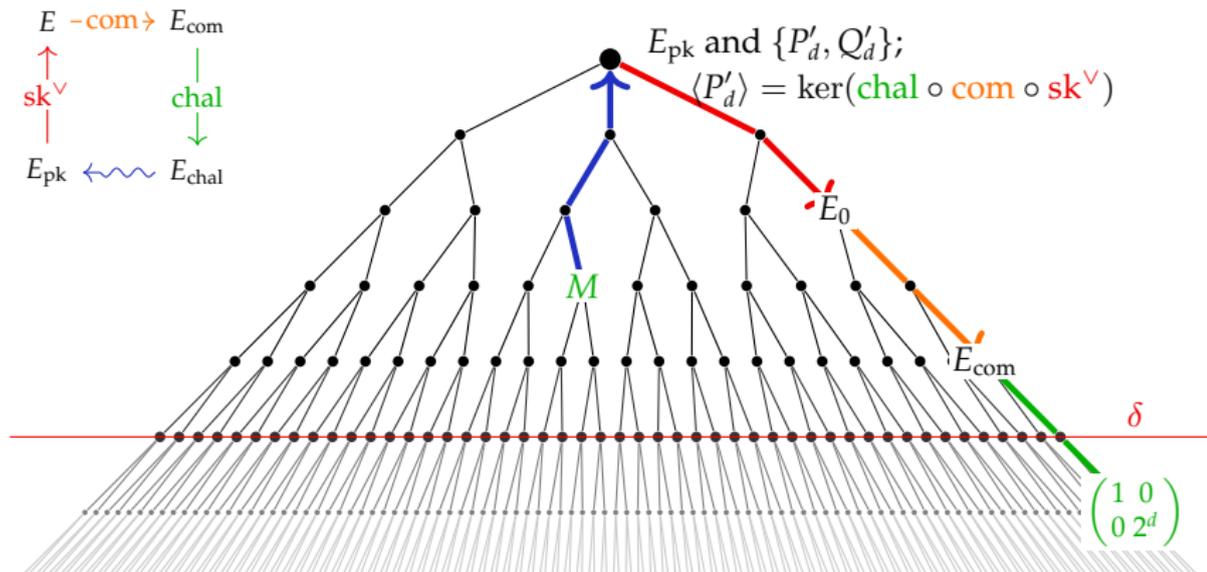
$$\langle P'_d \rangle = \ker(\text{chal} \circ \text{esk} \circ \text{sk}^\vee).$$

2. Using M , read off a shorter path back to the root, reveal the corresponding 2-isogeny chain.

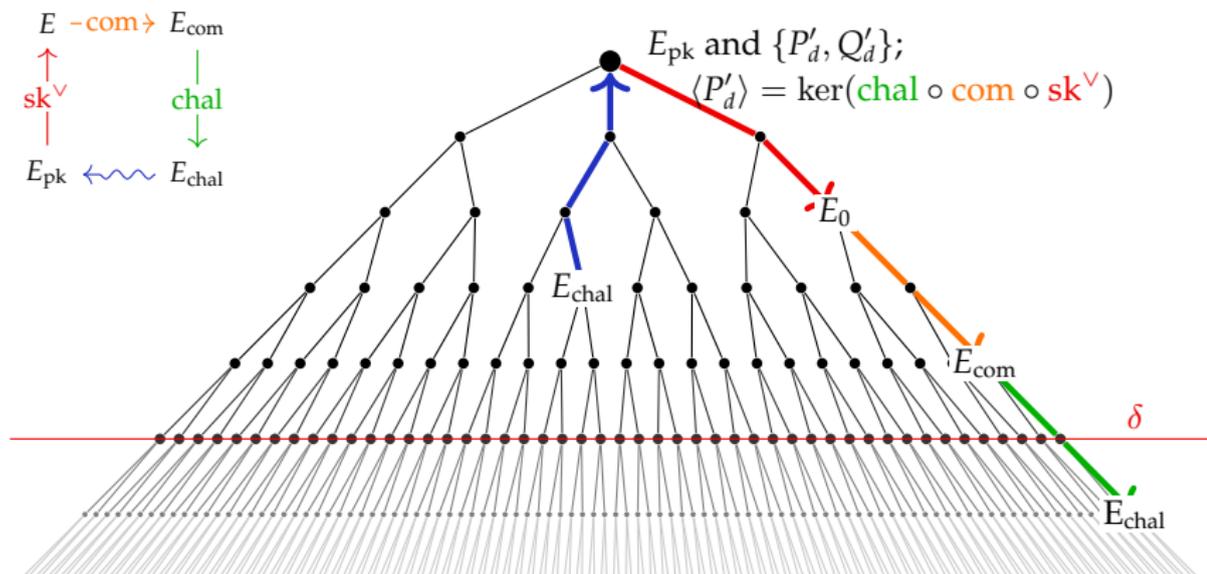
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In progress: computing quotients

- ▶ We can't yet compute $M \sim \begin{pmatrix} 1 & 0 \\ 0 & 2^d \end{pmatrix}$.
- ▶ Only existing algorithm: Franc-Masdeu.
- ▶ Adapting F-M leads to:

Open question:

Find the shortest vector of

$$\begin{pmatrix} N \\ 2^n I \end{pmatrix} \in \text{Mat}_{4 \times 8}(\mathbb{Z}[x])$$

with respect to length

$$\|\underline{v}(x)\| = \min_{x \in [0, r]} \{ \|\underline{v}(x) \cdot \underline{v}(x)\|_\infty \}.$$

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Questions? Answers?

References

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